

Social Welfare-based Optimization for Data/Service Delivery to Connected Vehicles via Edges

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and

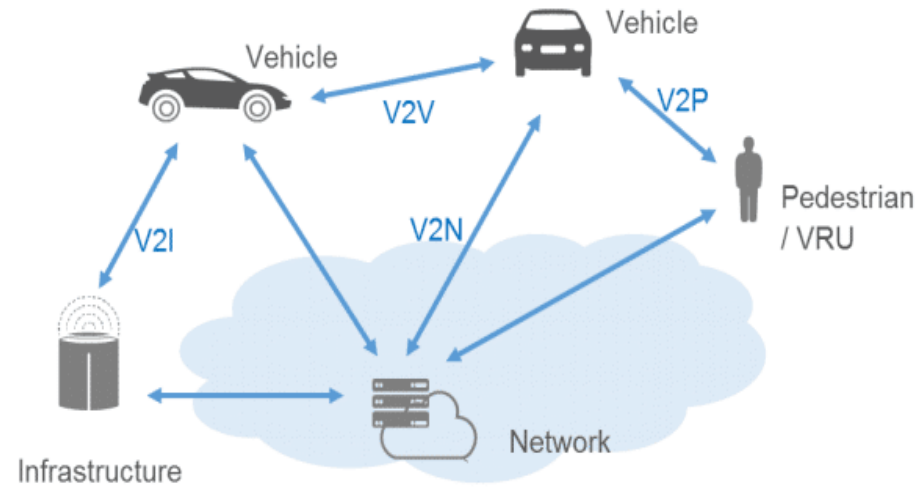
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Motivation: Increased Connectivity For Automotives

- Vehicle connectivity with various entities
 - to other vehicles → V2V
 - to infrastructure → V2I
 - to cloud → V2N/V2C
 - to pedestrian → V2P

- Potential for efficient delivery of data/services!



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Motivation: U.S. Government Policy Trend

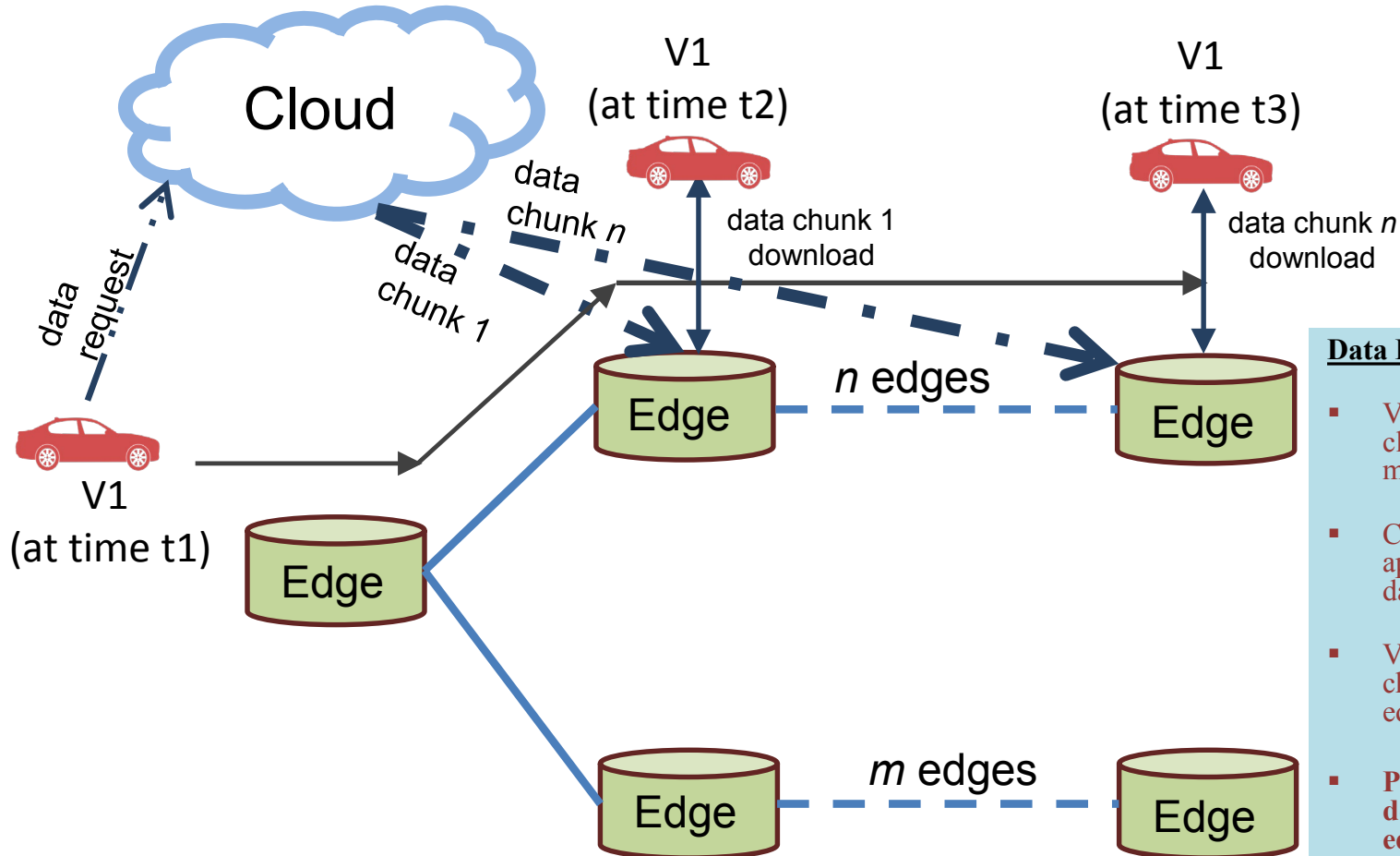
- USDOT *Connected* Vehicle Pilot Projects (~\$42 million)
 - New York City DOT Pilot
 - Tampa-Hillsborough Expressway Authority Pilot
 - Wyoming DOT Pilot
- New York City DOT Pilot
 - **(Goal)** Eliminate traffic related deaths and reduce crash
 - **(Approach)** Exchange information about hazardous situations using Vehicle-to-Vehicle (V2V), Vehicle-to-Infrastructure (V2I), Infrastructure-to-Pedestrian (IVP) communications

*Approx. 300 road side units (RSU) will be installed in NY city



*Source: www.cvp.nyc

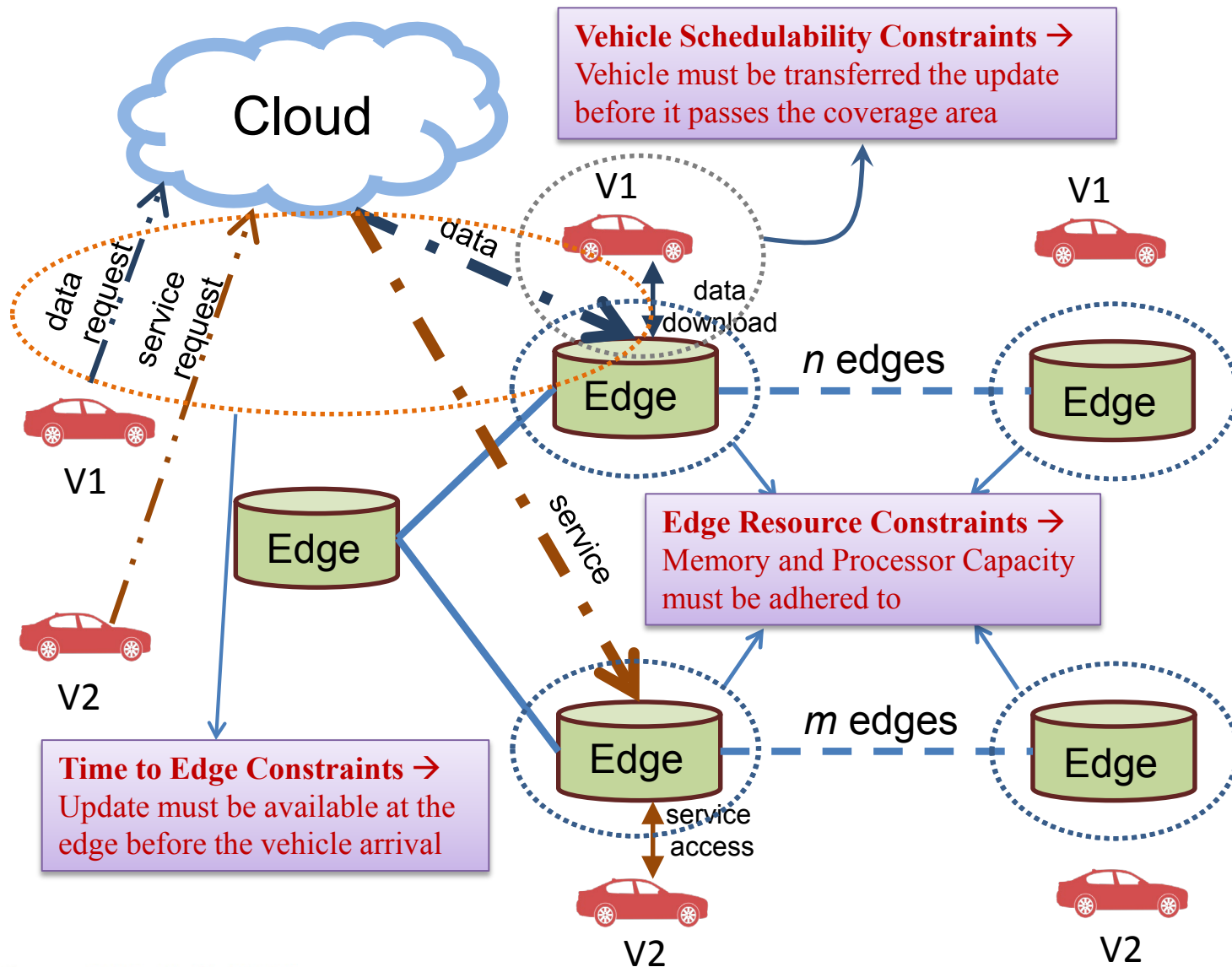
Motivation: Data Delivery Scenario



Data Delivery

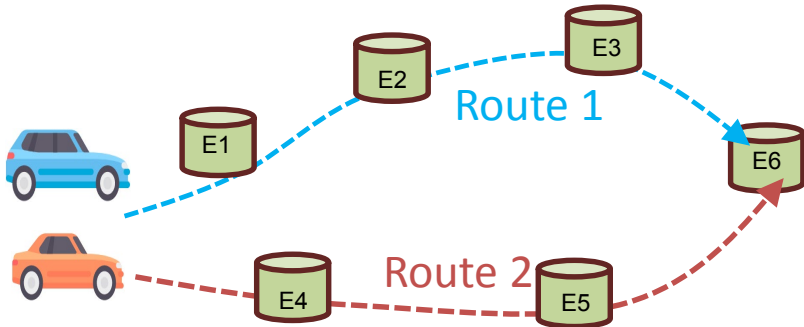
- Vehicle V1 requests the cloud for some data such as map data
- Cloud determines the appropriate edge to send the data
- V1 downloads the data chunks from the respective edges as it passes by them
- **Problem: Necessary to determine the appropriate edges to push the update chunks in order to optimize a user/system objective**

Problem Scenario – Potential Constraints

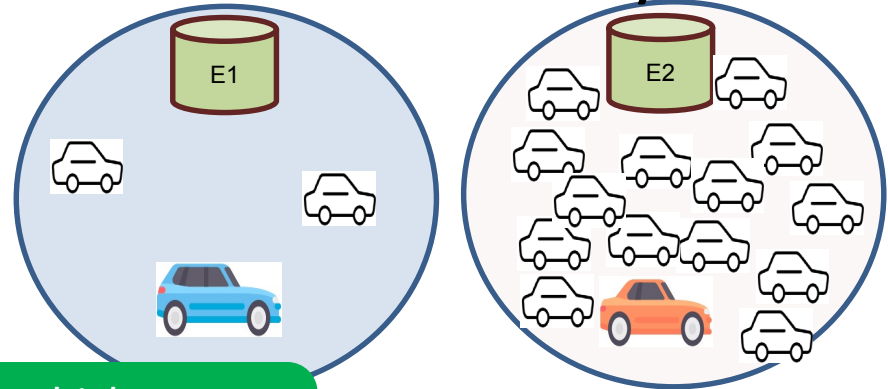


Problem Scenario – Influencing Factors

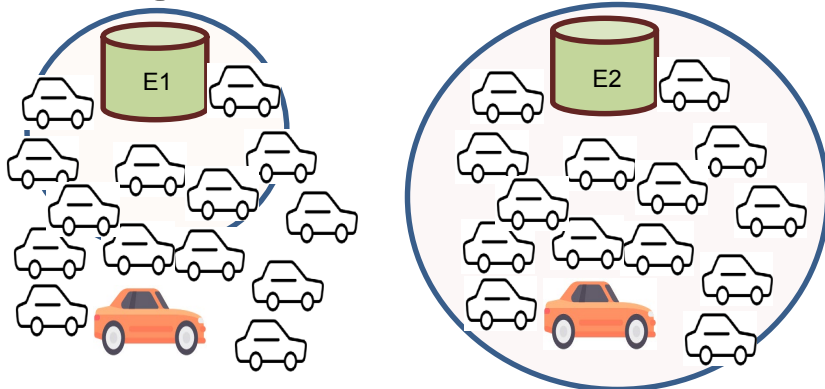
Vehicle routes



Vehicle density

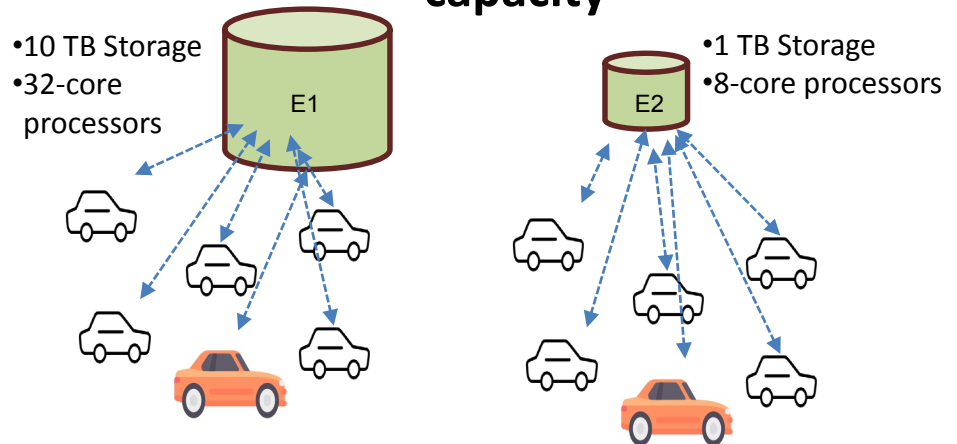


Field communication coverage



Can a group of vehicles successfully get data/services on their routes?

Memory/computing capacity



Optimization Problem Formulation – Constraints for Data/Service Delivery

- Problem size parameters

N – Number of vehicles, M – Number of edges

- Route indicator

$x_{i,j}$ – Binary value indicating whether vehicle i passes through edge j

(1 indicates that it passes and 0 indicates that it doesn't)

- Decision Variable

$m_{i,j}$ – Memory chunk of vehicle i allocated to edge j

$serv_{i,j}$ – Binary decision variable denoting vehicle i receives its service from edge j

Optimization Problem Formulation – Constraints for Data/Service Delivery

- Range Constraint

m_{app_i} – Memory requirement for vehicle i update

$$\left. \begin{array}{l} m_{i,j} \geq 0 \\ m_{i,j} \leq m_{app_i} \end{array} \right\} \forall i,j: x_{i,j} = 1$$

Memory chunk allocated must take a value between 0 and m_{app_i} both values included

- Edge Memory Constraint (Update Delivery)

m_{occ_j} - Size of memory currently occupied on edge j

m_{e_j} - Total memory size of edge j

$$\sum_{i=1}^N x_{i,j} * m_{i,j} + m_{occ_j} \leq m_{e_j}, \forall j = 1 \text{ to } M$$

Sum of all memory chunks for all vehicles on edge j cannot exceed the available memory

Optimization Problem Formulation – Constraints for Data/Service Delivery

- Edge Memory Constraint (Service Delivery)

d_i - Size of input data from vehicle to edge

r_i - Size of output data from edge to vehicle

$$\sum_{i=1}^N serv_{i,j} * (d_i + r_i) + m_{occ_j} \leq m_{e_j} , \forall j = 1 \text{ to } M$$

Sum of memory required for input and output data for all vehicles on edge j cannot exceed the available memory

- Edge Computation Constraint (Service Delivery)

P_{occ_j} - Processing resources already in use on edge j (can be number of processing units in time or number of VMs available)

P_j - Total processing capacity of edge j

p_i - Processing resource requirement of vehicle i

$$\sum_{i=1}^N serv_{i,j} * p_i + P_{occ_j} \leq P_j , \forall j = 1 \text{ to } M$$

Optimization Problem Formulation – Constraints for Data/Service Delivery

- Accumulation Constraint

$$\sum_{j=1}^M x_{i,j} * m_{i,j} = m_{app_i}, \forall i = 1 \text{ to } N$$

Sum of memory chunks for a vehicle over all edges should be equal to the update size

- Time to Edge Constraint

$t_{comm_{i,j}}$ – Time required to send update data to the edge j after vehicle i initiates update download

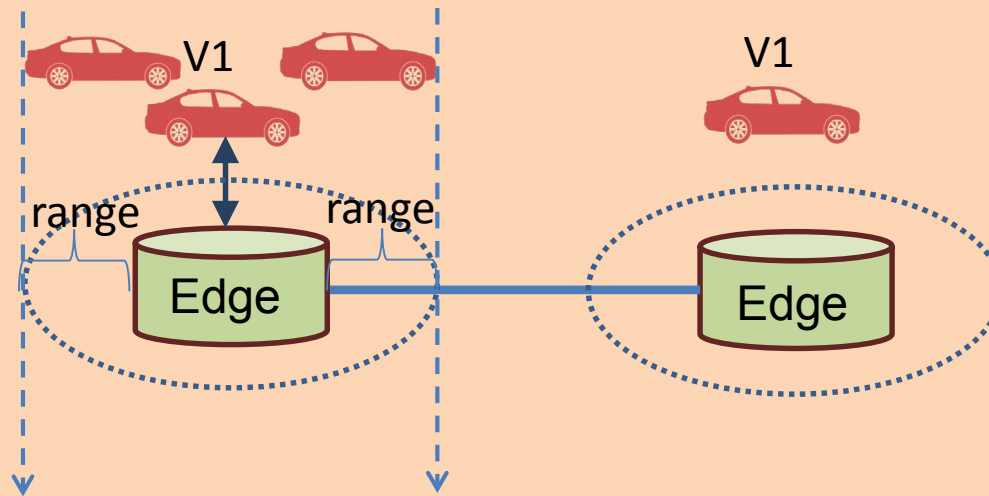
$t_{trv_{i,j}}$ – Time required to travel to edge j by vehicle i after initiating the download

$$m_{i,j} * t_{comm_{i,j}} \leq m_{i,j} * t_{trv_{i,j}}, \forall i, j$$

Update chunks must be available at the edges before the vehicle arrives at the edges

Optimization Problem Formulation – Constraints for Data/Service Delivery

- Bandwidth Schedulability Constraint



Macroscopic Flow Model

k – vehicle density, i.e., number of vehicles per unit distance
 q – vehicle flow, i.e., number of vehicles passing a fixed point per unit time
 v – Vehicle speed

$$q = k \times v$$

Optimization Problem Formulation – Constraints for Data/Service Delivery

- Bandwidth Schedulability Constraint

- Speed Density Relationship

$$v = v_f \times \left(1 - \frac{k}{k_{jam}}\right),$$

where v_f is the free flow speed (or speed limit) and k_{jam} is the vehicle density during a traffic jam

- Vehicles in the coverage area can be modelled as a queuing network, specifically in this scenario as a M/D/C queue (poisson arrival, deterministic service and C servers)
- As per M/D/C queue modelling, the minimum number of bytes received by a vehicle $D_{min} = B / (k_{jam} \times v)$ [1], where B is the bandwidth of the edge device
- **The bandwidth schedulability constraint for vehicle requiring update is**

$$m_{i,j} \leq D_{min}$$

[1] Analytical Models and Performance Evaluation of Drive-thru Internet Systems, IEEE JSAC 2011

Optimization Problem Formulation – Constraints for Data/Service Delivery

- Bandwidth Schedulability Constraint

- For service delivery, as per M/D/C modelling, the minimum number of bytes received by a vehicle is

$$D_{min}^{serv} = \frac{B \times \left(\frac{L}{v} - t_{p_i} \right)}{k_{jam} \times L}$$

where L is the coverage distance, t_{p_i} is the time taken to execute a service for vehicle i with p_i number of VMs

- **The bandwidth schedulability constraint for vehicle requiring a service is given by**

$$serv_{i,j} * (d_i + r_i) \leq D_{min}^{serv}$$

Edge Bandwidth Utilization

- Total bandwidth utilization can then be calculated as

$$bw^{util} = bw^{util,data} + bw^{util,serv} + bw^{util,oth}$$

where

$bw^{util,data}$ - Bandwidth utilization of the vehicles requiring data delivery

$bw^{util,serv}$ - Bandwidth utilization of the vehicles requiring service delivery

$bw^{util,oth}$ - Bandwidth utilization of the vehicles that do not require service or data delivery

Optimization Objective

- Congestion Cost Objective

Minimization of total bandwidth cost considering all edges

- Motivation : In order to avoid congestion of vehicles at an edge, a cost must be determined in utilizing the bandwidth resource given by $bw_j^{cost} = \beta(1 + bw_j^{util})^2$ [2], where bw_j^{util} is the bandwidth utilization on edge j and β is the bandwidth cost factor
- **The optimization objective is**

$$\text{minimize } \sum_{j=1}^M bw_j^{cost}$$

[2] On-Demand Bandwidth Pricing for Congestion Control in Core Switches in Cloud Networks, IEEE CLOUD 2016

Social Welfare Function for Data/Service Delivery in Connected Vehicles

- Two functions used to compute social welfare – **Utility function** and **Congestion Cost function**
- **Utility Function**
 - A function of the data allocated for a vehicle on the edge given by $U_{i,j}(m_{i,j})$
 - A concave function is used such that the utility increases as the allocated data increases
 - $\log(1 + m_{i,j})$ is the preferred choice based on earlier works
- In order to capture the effect of faster delivery, we must use different utility functions at different edges for a vehicle

Reason – The utility value must decrease across edges on a route for vehicles which need faster delivery → So $U_{i,j}(m_{i,j})$ scales down as we go from one edge to the next

Social Welfare Function

- **Congestion Cost Function** – Same as the bandwidth congestion cost function used earlier denoted by $C_j(bw_{util,j}) \rightarrow$ Is a convex function
- Social Welfare Function

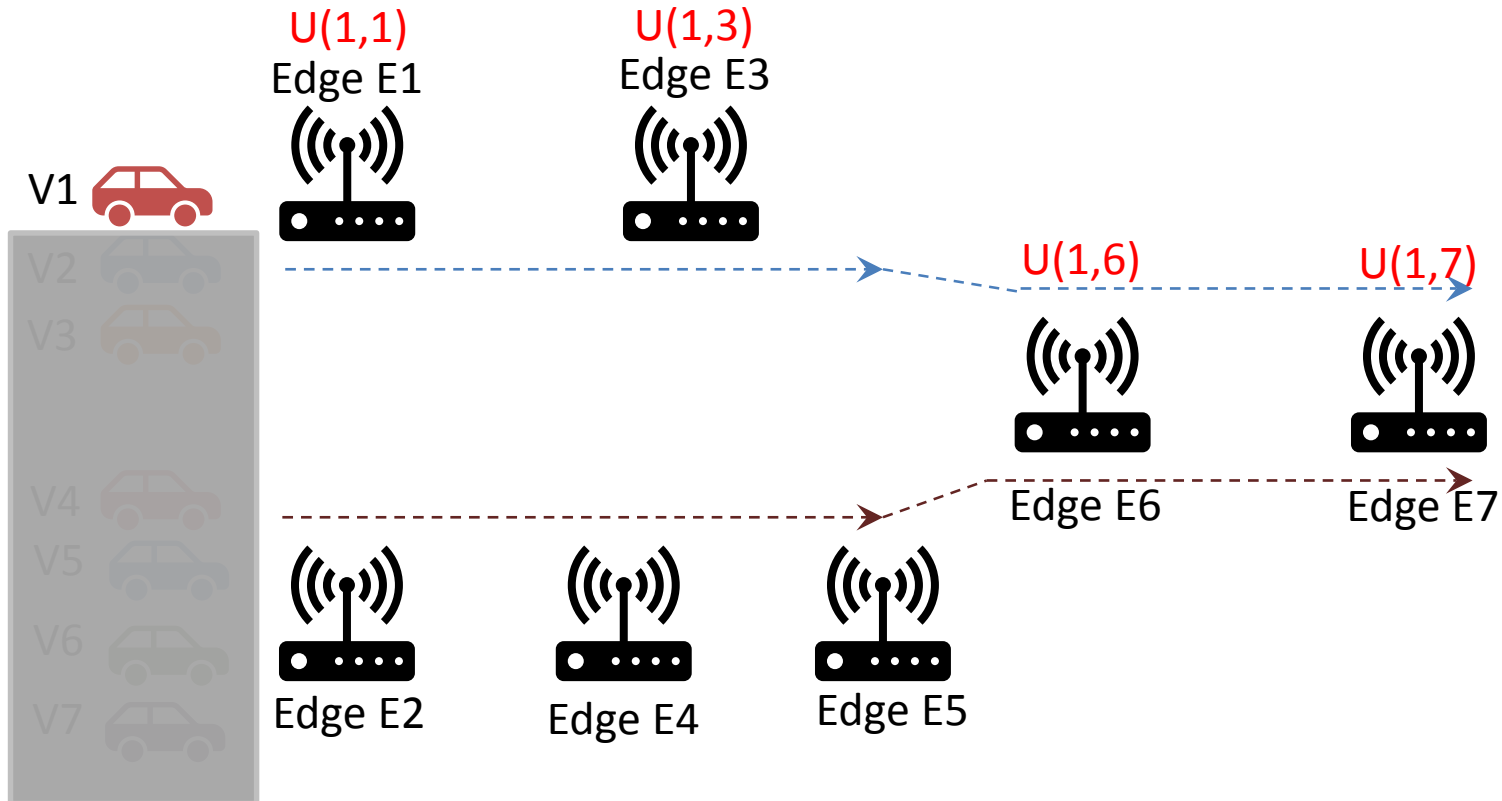
$$S(B) = \sum_{i \in V} \sum_{j \in E} U_{i,j}(m_{i,j}) - \sum_{j \in E} C_j(bw_{util,j})$$

where $U_{i,j}(m_{i,j}) = \log(1 + m_{i,j}) \times 1/(s \times t_{trv_{i,j}})$ and s is the scaling factor determining the utility of the edge

- It is a concave function and the objective function will be to **maximize $S(B)$**

What does it mean?

Utility Values



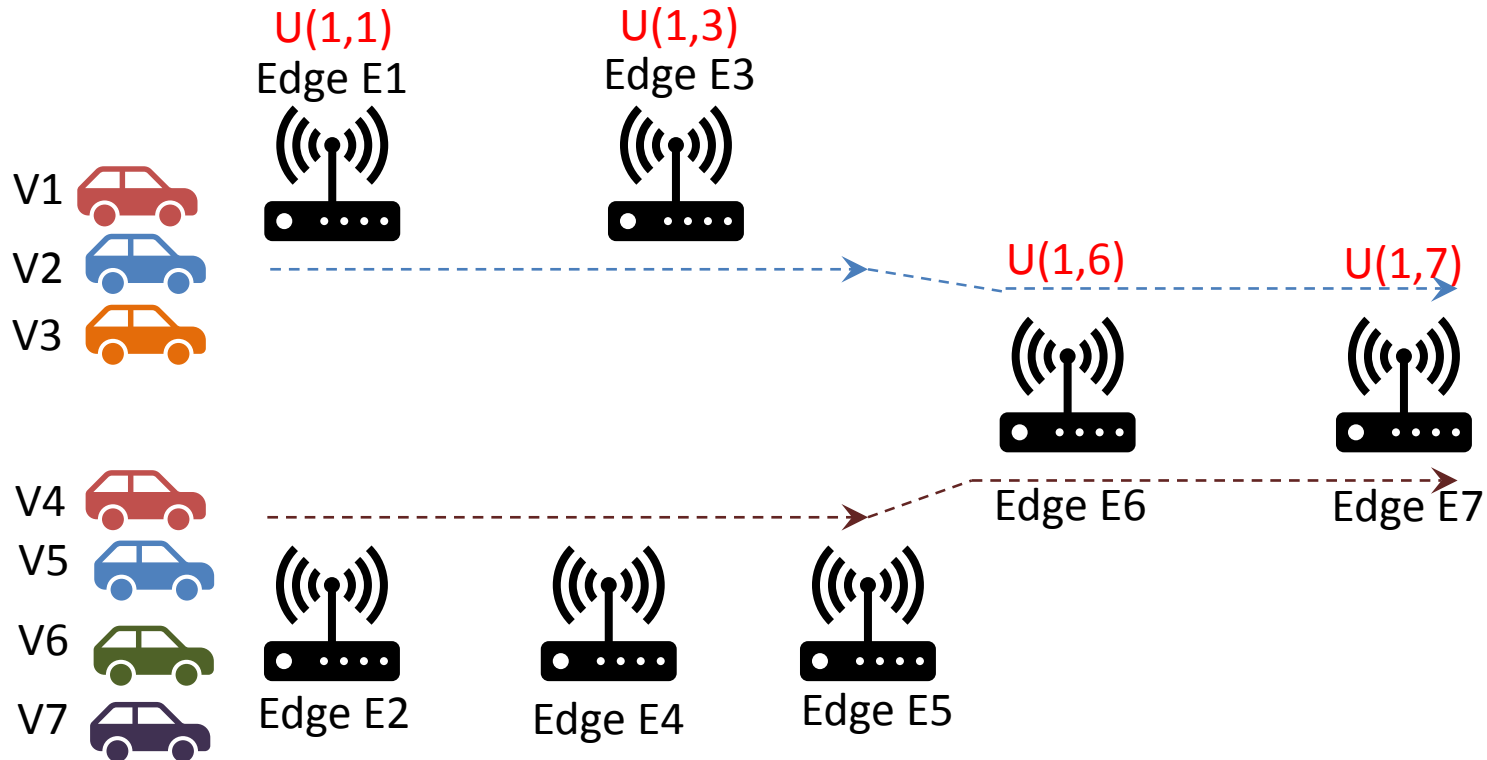
Utility value $U(1,1) = \frac{1}{t(1,1)} * \log\{1 + m(1,1)\}$, where

$t(1,1)$ - earliest time when V1 will reach E1, and

$m(1,1)$ - Data allocated to V1 on E1

What does it mean?

Utility Values



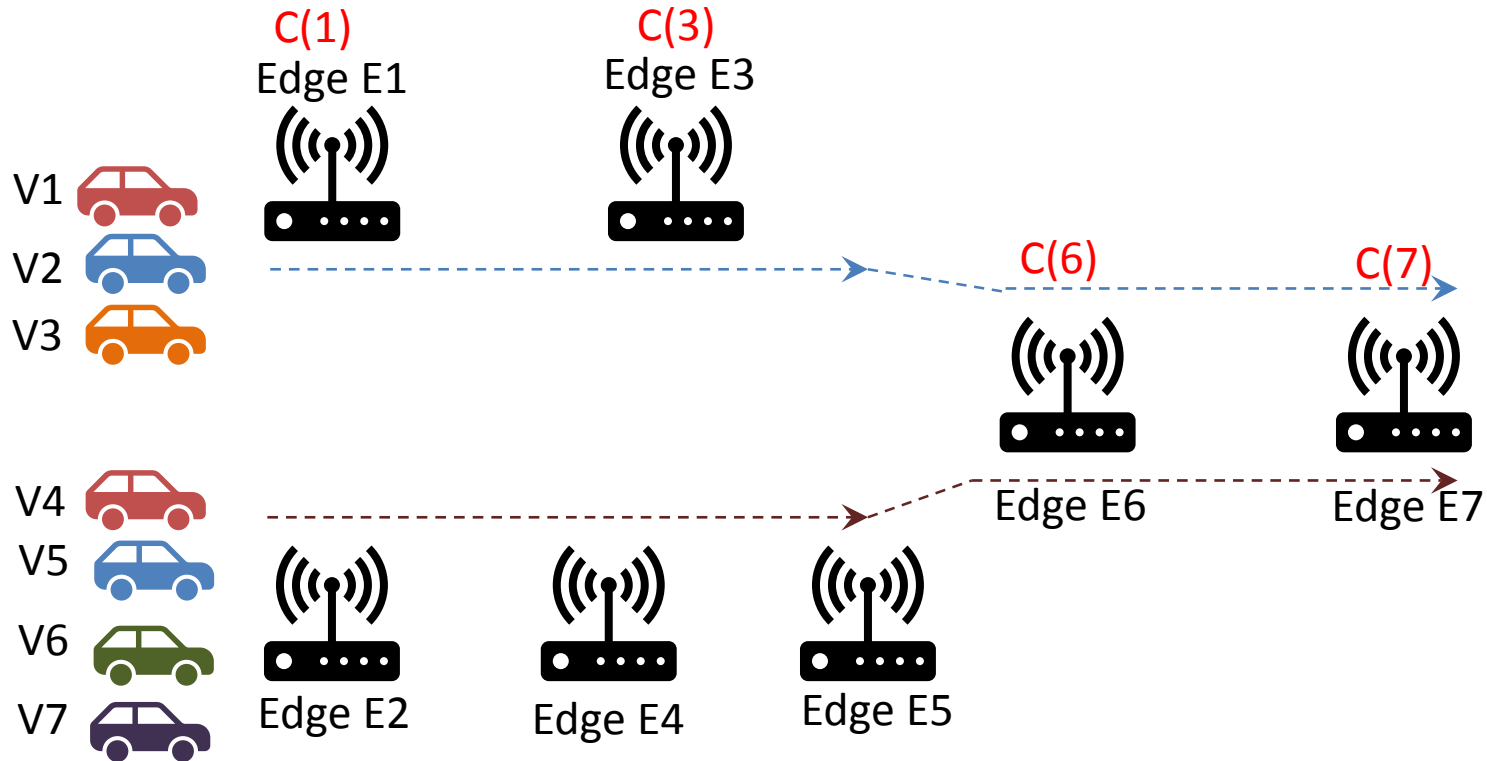
Total Utility for $TU(V1) = U(1,1) + U(1,3) + U(1,6) + U(1,7)$

Similarly total utility is calculated for each vehicle.

$$\text{System Utility} = \sum_{\forall V_i \in \text{vehicles}} TU(V_i)$$

What does it mean?

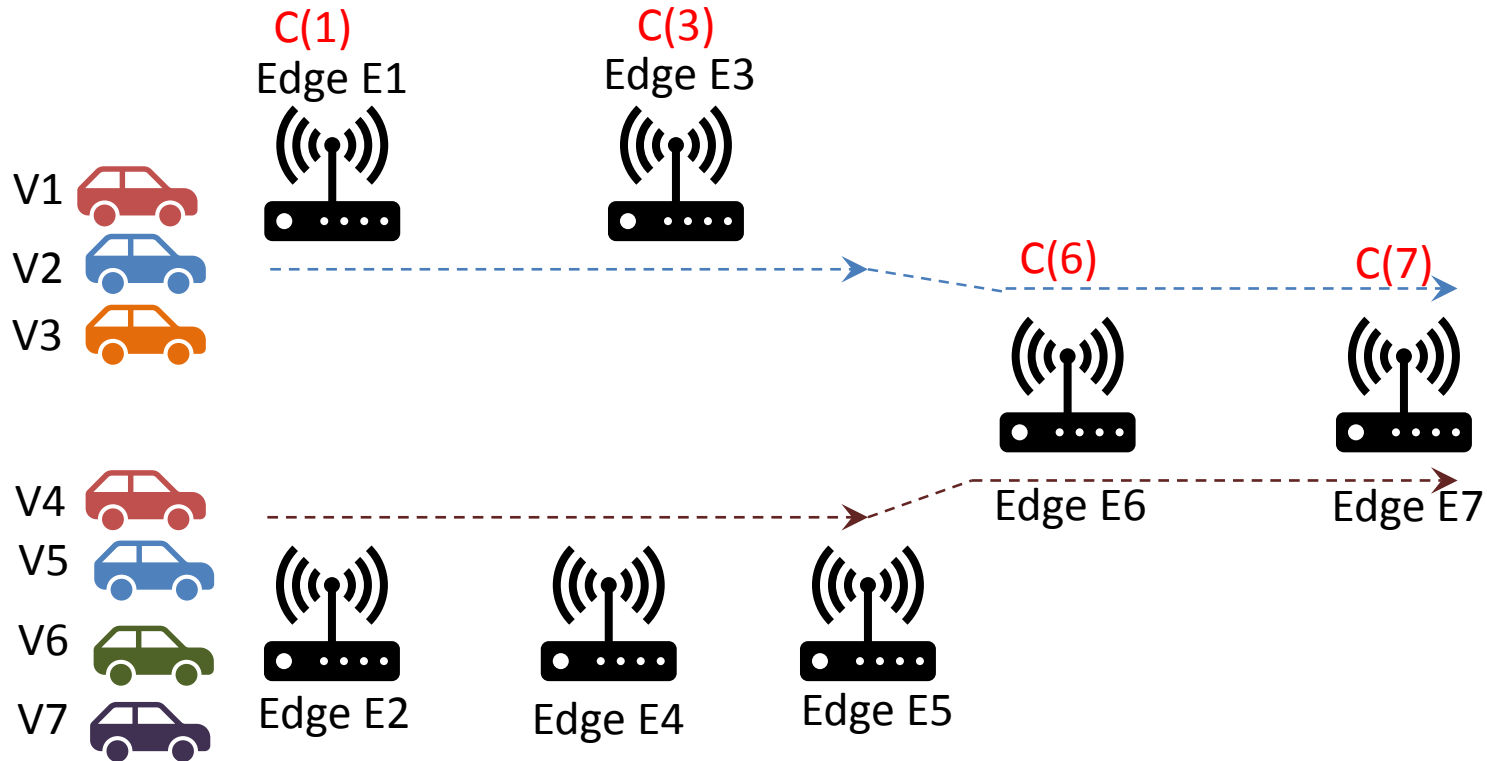
Congestion Cost



Congestion Cost on Edge E1 $C(1) = \beta * (1 + BU(1))^2$, where
 β – Congestion cost factor, and
 $BU(1)$ – bandwidth utilization due to V1, V2 and V3 on E1

What does it mean?

Congestion Cost



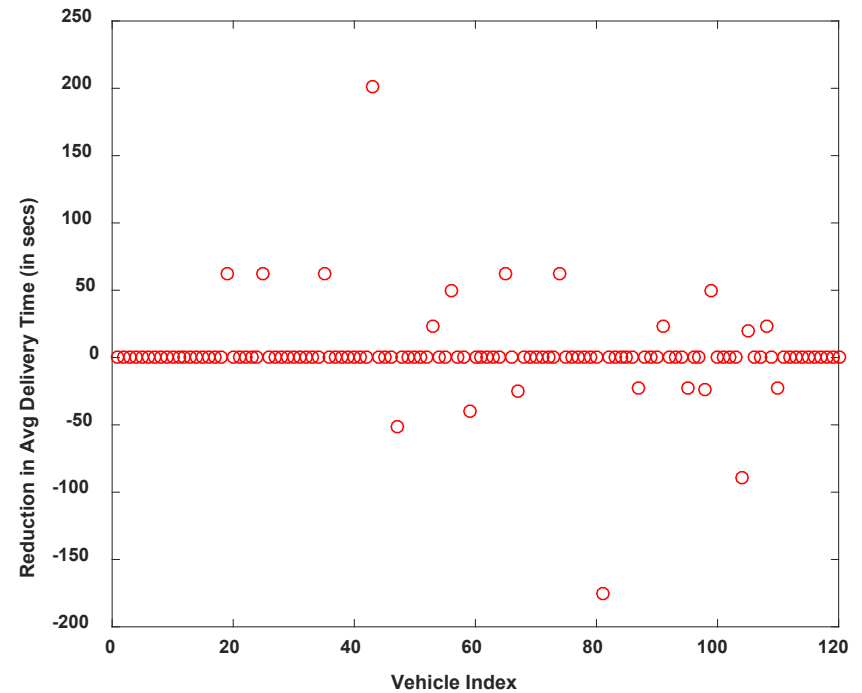
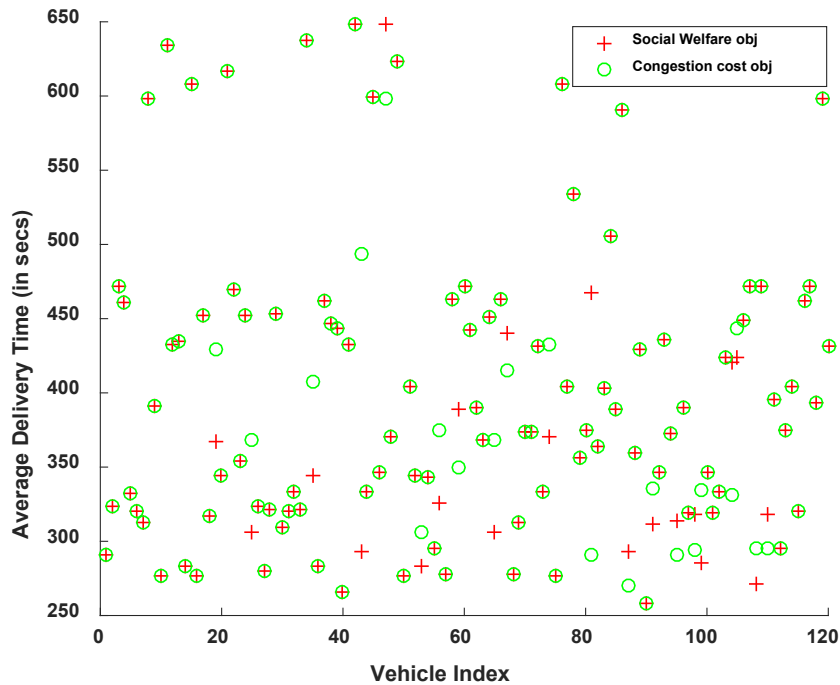
$$\text{Total System Congestion Cost} = \sum_{j \in \text{edges}} C(j)$$

$$\text{Social Welfare Function} = \text{Total System Utility} - \text{Total System Congestion Cost}$$

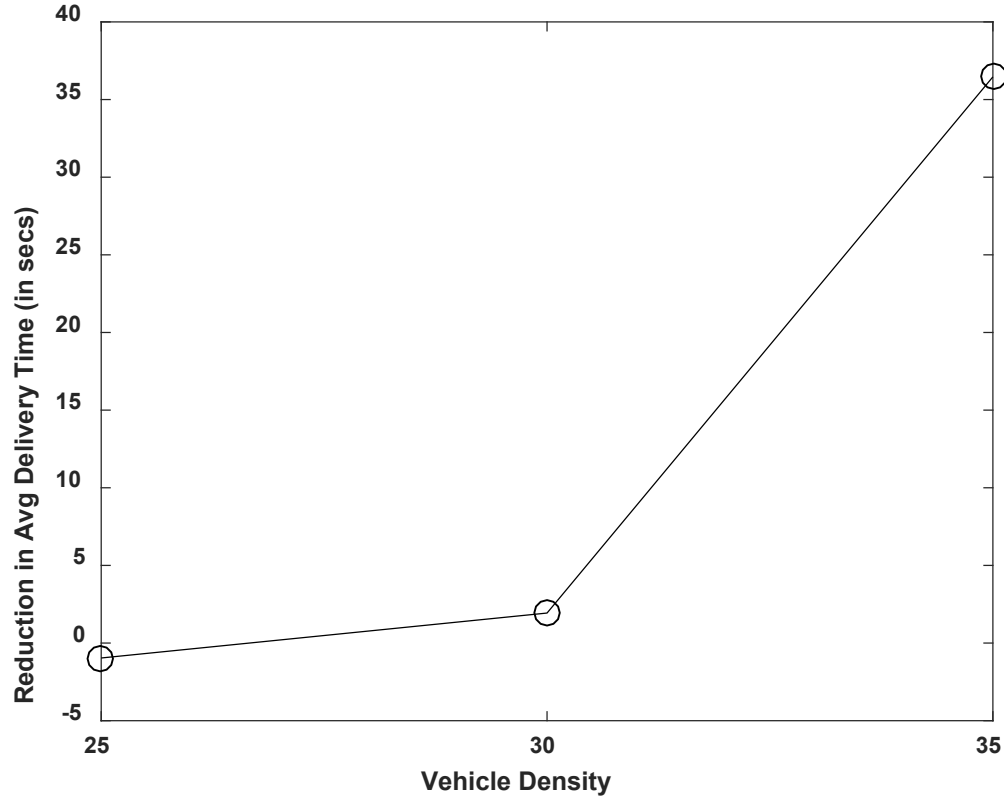
Experimental Results

- Experimental Setting 1
 - Number of edges (M) = {25,49,81,100}
 - Number of vehicles (N) requiring update was varied between 20 to 200
 - Vehicle jam density (k_{jam}) was randomly generated between 40 and 50 at all edges
 - The actual vehicle density in the coverage area was varied between 35 and 25 in steps of 5
 - Coverage distance (L_j) was assigned arbitrary values between 0.6 miles and 1.6 miles
 - Memory requirement of data requested by each vehicle was randomly generated between 60 and 80 Mbits
 - Memory capacity of the edge was randomly generated between 400 and 500 Mbits
 - Maximum bandwidth capacity of the edges was randomly generated between 8 and 15 Mbps
 - Free flowing velocity of the vehicles at the edges was randomly generated between 50 and 70 mph
 - The route of the vehicles was randomly generated by picking connected edges in a grid

Result 1: Data retrieval time using two objective functions



Result 2: Variation in Delivery Time Reduction with Vehicle Densities



- At lower vehicle densities, the average delivery time considering all vehicles is marginally higher
- As vehicle densities increase, the social welfare objective results in higher reduction in average delivery time
- For vehicle density = 35, the reduction is 36.46 secs

Conclusion

- Proposed an optimization framework for data/service delivery to connected vehicles considering vehicle flow densities
- Proposed a social welfare based delivery
- Demonstrated significant reductions in average delivery times with increasing vehicle densities

Thank You!
Questions

